

# Applied Physics 158

## Numerical Linear Algebra

1<sup>st</sup> sem, A.Y. 2020–2021

## 1 Methods

Linear systems can generally be written in the form

$$Ax = b \quad (1)$$

where  $A$  is an  $m \times n$  matrix,  $x$  is a vector of length  $n$ , and  $b$  is a vector of length  $m$ . There can be no solution, one solution, or infinitely many solutions for such configuration. In matrix form this can be written as

$$\begin{pmatrix} a_{1,1} & a_{1,2} & \cdots & a_{1,n-1} & a_{1,n} \\ a_{2,1} & a_{2,2} & \cdots & a_{2,n-1} & a_{2,n} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{m-1,1} & a_{m-1,2} & \cdots & a_{m-1,n-1} & a_{m-1,n} \\ a_{m,1} & a_{m,2} & \cdots & a_{m,n-1} & a_{m,n} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_{n-1} \\ x_n \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_{m-1} \\ b_m \end{pmatrix} \quad (2)$$

which corresponds to the linear system

$$\begin{aligned} a_{1,1}x_1 + a_{1,2}x_2 + \cdots + a_{1,n-1}x_{n-1} + a_{1,n}x_n &= b_1 \\ a_{2,1}x_1 + a_{2,2}x_2 + \cdots + a_{2,n-1}x_{n-1} + a_{2,n}x_n &= b_2 \\ &\vdots \\ a_{m-1,1}x_1 + a_{m-1,2}x_2 + \cdots + a_{m-1,n-1}x_{n-1} + a_{m-1,n}x_n &= b_{m-1} \\ a_{m,1}x_1 + a_{m,2}x_2 + \cdots + a_{m,n-1}x_{n-1} + a_{m,n}x_n &= b_m \end{aligned} \quad (3)$$

### 1.1 Gaussian Elimination

For simplicity, let  $A$  be a square matrix. The premise of Gaussian elimination is to determine from the augmented matrix

$$\left( \begin{array}{ccccc|c} a_{1,1} & a_{1,2} & \cdots & a_{1,n-1} & a_{1,n} & b_1 \\ a_{2,1} & a_{2,2} & \cdots & a_{2,n-1} & a_{2,n} & b_2 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ a_{n-1,1} & a_{n-1,2} & \cdots & a_{n-1,n-1} & a_{n-1,n} & b_{n-1} \\ a_{n,1} & a_{n,2} & \cdots & a_{n,n-1} & a_{n,n} & b_n \end{array} \right) \quad (4)$$

a row echelon form that can easily give the solutions of the system of equations. This is shown by

$$\left( \begin{array}{ccccc|c} a'_{1,1} & a'_{1,2} & \cdots & a'_{1,n-1} & a'_{1,n} & b'_1 \\ 0 & a'_{2,2} & \cdots & a'_{2,n-1} & a'_{2,n} & b'_2 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & a'_{n-1,n-1} & a'_{n-1,n} & b'_{n-1} \\ 0 & 0 & \cdots & 0 & a'_{n,n} & b'_n \end{array} \right) \quad (5)$$

and can be calculated by following the rules of solving linear equations:

- a nonzero constant may be multiplied on both sides of any of the equation,
- two rows can be switched with their corresponding  $b$  values, and
- adding and/or subtracting two rows are allowed and placing it on one of those rows.

An algorithmic way to solve this is by looping the successive assignment

$$a_{i,j} = a_{i,j} - \frac{a_{i,k}}{a_{k,k}} a_{k,j} \quad (6)$$

for each step  $k$  from  $k = 1, 2, \dots, n-1$  for values  $i = k+1, \dots, n$  and  $j = k+1, \dots, n+1$ .

This can be decomposed further into a reduced row echelon form

$$\left( \begin{array}{ccccc|c} a''_{1,1} & 0 & \cdots & 0 & 0 & b''_1 \\ 0 & a''_{2,2} & \cdots & 0 & 0 & b''_2 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & a''_{n-1,n-1} & 0 & b''_{n-1} \\ 0 & 0 & \cdots & 0 & a''_{n,n} & b''_n \end{array} \right) \quad (7)$$

which creates a diagonal matrix that easily gives away the value of each  $x$  element. This is a variant of the Gaussian elimination called the Gauss-Jordan elimination. Algorithmically, this is given by

$$a_{i,j} = \begin{cases} a_{i,n+1} - \frac{a_{i,k}}{a_{k,k}} a_{k,n+1} & i = 1, 2, \dots, k-1, k+1, \dots, n \\ a_{i,n+1} & i = k \end{cases} \quad (8)$$

for each step  $k$  from  $k = 1, 2, \dots, n-1$ .

## 1.2 LU Decomposition

Suppose the square matrix  $A$  can be broken down into two factors  $L$  and  $U$  such that

$$A = LU \quad (9)$$

and the decomposition is composed of a lower triangular matrix  $L$  and an upper triangular matrix  $U$ . This technique is called an LU decomposition or the Doolittle's method. The matrix multiplication  $LU$  can be written as

$$\begin{pmatrix} 1 & 0 & \cdots & 0 & 0 \\ l_{2,1} & 1 & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ l_{n-1,1} & l_{n-1,2} & \cdots & 1 & 0 \\ l_{n,1} & l_{n,2} & \cdots & l_{n,n-1} & 1 \end{pmatrix} \begin{pmatrix} u_{1,1} & u_{1,2} & \cdots & u_{1,n-1} & u_{1,n} \\ 0 & u_{2,2} & \cdots & u_{2,n-1} & u_{2,n} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & u_{n-1,n-1} & u_{n-1,n} \\ 0 & 0 & \cdots & 0 & u_{n,n} \end{pmatrix} \quad (10)$$

This can be simplified into the summation notation

$$a_{i,j} = \sum_{k=1}^n l_{i,k} u_{k,j} \quad (11)$$

As an algorithm, the values can be determined by the row assignment for  $r$  in  $r = 1, 2, \dots, n$  of  $U$

$$u_{r,j} = a_{r,j} - \sum_{k=1}^{r-1} l_{r,k} u_{k,j} \quad (12)$$

for  $j = r, r+1, \dots, n$ . The columns of  $L$  can immediately be assigned for  $r$  in  $r = 1, 2, \dots, n-1$

$$l_{i,r} = \frac{a_{i,r} - \sum_{k=1}^{r-1} l_{i,k} u_{k,r}}{u_{r,r}} \quad (13)$$

for  $i = r+1, r+2, \dots, n$ .

## 1.3 Cholesky and LDL Decomposition

Suppose the square matrix  $A$  can be broken down into two factors  $L$  and  $L^T$  such that

$$A = LL^T \quad (14)$$

and the decomposition is composed of a lower triangular matrix  $L$  and its transpose  $L^T$ . This technique is called the Cholesky Decomposition. The matrix multiplication  $LL^T$  can be written as

$$\begin{pmatrix} l_{1,1} & 0 & \cdots & 0 & 0 \\ l_{2,1} & l_{2,2} & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ l_{n-1,1} & l_{n-1,2} & \cdots & l_{n-1,n-1} & 0 \\ l_{n,1} & l_{n,2} & \cdots & l_{n,n-1} & l_{n,n} \end{pmatrix} \begin{pmatrix} l_{1,1} & l_{2,1} & \cdots & l_{n-1,1} & l_{n,1} \\ 0 & l_{2,2} & \cdots & l_{n-1,2} & l_{n,2} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & l_{n-1,n-1} & l_{n,n-1} \\ 0 & 0 & \cdots & 0 & l_{n,n} \end{pmatrix} \quad (15)$$

By evaluating this matrix multiplication and comparing the elements to  $A$ , an algorithm can be created to determine the values of  $L$ .

Generally, the values on the diagonal of  $L$  can be determined by

$$l_{k,k} = \sqrt{a_{k,k} - \sum_{j=1}^{k-1} l_{k,j}^2} \quad (16)$$

for  $k$  in  $k = 1, 2, \dots, j$  and for the remaining elements below the diagonal of  $L$ ,

$$l_{i,k} = \frac{1}{l_{k,k}} \left( a_{i,k} - \sum_{j=1}^{k-1} l_{i,j} l_{k,j} \right) \quad (17)$$

for  $k$  in  $k = 1, 2, \dots, j$ , for  $j$  in  $j = 1, 2, \dots, i$ , and for  $i$  in  $i = 1, 2, \dots, n$ .

Related to this is the  $LDL$  decomposition. Suppose the square matrix  $A$  can be broken down into three factors  $L$ ,  $D$ , and  $L^T$  such that

$$A = LDL^T \quad (18)$$

and the decomposition is composed of a lower triangular matrix  $L$ , its transpose  $L^T$ , and a diagonal matrix  $D$ . The matrix multiplication  $LDL^T$  can be written as

$$\begin{pmatrix} 1 & 0 & \cdots & 0 & 0 \\ l_{2,1} & 1 & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ l_{n-1,1} & l_{n-1,2} & \cdots & 1 & 0 \\ l_{n,1} & l_{n,2} & \cdots & l_{n,n-1} & 1 \end{pmatrix} \begin{pmatrix} d_{1,1} & 0 & \cdots & 0 & 0 \\ 0 & d_{2,2} & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & d_{n-1,n-1} & 0 \\ 0 & 0 & \cdots & 0 & d_{n,n} \end{pmatrix} \begin{pmatrix} 1 & l_{2,1} & \cdots & l_{n-1,1} & l_{n,1} \\ 0 & 1 & \cdots & l_{n-1,2} & l_{n,2} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & 1 & l_{n,n-1} \\ 0 & 0 & \cdots & 0 & 1 \end{pmatrix} \quad (19)$$

Similar to Cholesky decomposition, by evaluating this matrix multiplication and comparing the elements to  $A$ , an algorithm can be created to determine the values of  $L$  and  $D$ .

## 1.4 Jacobi Method

Suppose the square matrix  $A$  can be broken down into three addends  $L$ ,  $D$ , and  $U$  such that

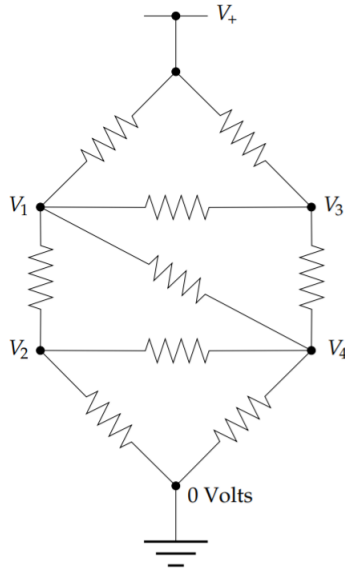
$$A = L + D + U \quad (20)$$

and the decomposition is composed of a lower triangular matrix  $L$ , a diagonal matrix  $D$ , and an upper triangular matrix  $U$ . This technique is called the Jacobi method. The matrix addition  $L + D + U$  can be written as

$$\begin{pmatrix} 0 & 0 & \cdots & 0 & 0 \\ a_{2,1} & 0 & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{n-1,1} & a_{n-1,2} & \cdots & 0 & 0 \\ a_{n,1} & a_{n,2} & \cdots & a_{n,n-1} & 0 \end{pmatrix} + \begin{pmatrix} a_{1,1} & 0 & \cdots & 0 & 0 \\ 0 & a_{2,2} & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & a_{n-1,n-1} & 0 \\ 0 & 0 & \cdots & 0 & a_{n,n} \end{pmatrix} + \begin{pmatrix} 0 & a_{1,2} & \cdots & a_{1,n-1} & a_{1,n} \\ 0 & 0 & \cdots & a_{2,n-1} & a_{2,n} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & 0 & a_{n-1,n} \\ 0 & 0 & \cdots & 0 & 0 \end{pmatrix} \quad (21)$$

## 2 Activity

Consider the following matrix:



### 3 Machine Problem – Circuit of Resistors

Consider the circuitry with resistors all with resistance  $R$  and a voltage  $V_+ = 5\text{ V}$ .

Using Kirchhoff's Junction Rule, the junction at  $V_1$  has the following equation

$$\frac{V_1 - V_2}{R} + \frac{V_1 - V_3}{R} + \frac{V_1 - V_4}{R} + \frac{V_1 - V_+}{R} = 0 \quad (22)$$

which is equivalent to the equation

$$4V_1 - V_2 - V_3 - V_4 = V_+ \quad (23)$$

Determine the three more equations for the junctions at  $V_2$ ,  $V_3$ , and  $V_4$ . The resulting system of equations can be used to solve for the values of the four voltages. Use any of the discussed methods to find these voltages.